

Here is Bubble Sort again:

```
In [1]: def bubble_sort(L):
        for i in range(len(L)):
            #Set swapped to False. If nothing changes during a pass, we break.
            #If that happens, the runtime is better than the worst-case runtime
            #for a list of size n
            swapped = False

            #in the worst case, we need n-1 passes
            #The worst case happens when the smallest element is at L[-1], since
            #the smallest element only shifts one position to the left with each pass
            for j in range(len(L)-1-i):
                if L[j] > L[j+1]:
                    L[j], L[j+1] = L[j+1], L[j]
                    swapped = True
            if not swapped:
                break
```

## Runtime Complexity Analysis -- Bubble Sort

Unlike with Selection Sort, Bubble Sort *can* terminate early -- if we break because a sweep didn't result in any two elements being swapped, the function returns faster.

We know that Bubble Sort will not run for more than  $n$  sweeps (where  $n = \text{len}(L)$ ), just because the outer loop will not run for more than  $n$  iterations. (Recall that in the previous lecture we argued that Bubble Sort only ever needs at most  $n-1$  sweeps.

Is it ever the case that Bubble Sort needs all  $n-1$  sweeps to sort the list? Yes.

Note what happens when the smallest element of the list is initially the last element of the list:

```

In [2]: def bubble_sort(L):
        for i in range(len(L)):
            swapped = False
            for j in range(len(L)-1-i):
                if L[j] > L[j+1]:
                    L[j], L[j+1] = L[j+1], L[j]
                    print("Swapped", L[j], "and", L[j+1])
                    print(L)
                    swapped = True
                else:
                    print("No need to swap", L[j], "and", L[j+1])
                    print(L)
            if not swapped:
                return

        print("=====")
if __name__ == '__main__':
    L = [2, 3, 4, 5, 1]
    bubble_sort(L)

```

```

No need to swap 2 and 3
[2, 3, 4, 5, 1]
No need to swap 3 and 4
[2, 3, 4, 5, 1]
No need to swap 4 and 5
[2, 3, 4, 5, 1]
Swapped 1 and 5
[2, 3, 4, 1, 5]
=====
No need to swap 2 and 3
[2, 3, 4, 1, 5]
No need to swap 3 and 4
[2, 3, 4, 1, 5]
Swapped 1 and 4
[2, 3, 1, 4, 5]
=====
No need to swap 2 and 3
[2, 3, 1, 4, 5]
Swapped 1 and 3
[2, 1, 3, 4, 5]
=====
Swapped 1 and 2
[1, 2, 3, 4, 5]
=====

```

Note that the 1 only moves left *once* per sweep. That makes sense: any element can only move left once during the sweep (but an element can move to the right many times.)

We can therefore conclude that in the **worst case**, Bubble Sort does not return before performing all  $n$  iterations of the outer loop.

Let's now figure out the worst-case runtime complexity of Bubble Sort for a list of length  $n$  by counting how many times the inner block repeats.

At iteration 0, the block runs for  $n-1-0$  times At iteration 1, the block runs for  $n-1-1$  times At iteration  $n-1$ , the block runs for  $n-1-(n-1)=0$  times

So in total, the block runs  $\sum_{i=0}^{n-1} (n - i - 1)$  times.

$$\sum_{i=0}^{n-1} (n - i - 1) = n^2 - \sum_{i=0}^{n-1} i - n = n^2 - n(n - 1)/2 - n = n^2 - n^2/2 + n/2 - n = n^2/2 - n/2$$

(To compute  $\sum_{i=0}^{n-1} i$ , we use the fact that  $\sum_{j=1}^m j = m(m + 1)/2$ , and substitute  $m = n - 1$ .)

We can therefore conclude that the worst-case runtime complexity of Bubble Sort is  $\mathcal{O}(n^2)$ , just like Selection Sort.